## Algorithm Closest_pair(PX, n)

//Input: An array PX[0 .. n-1] of points, and parameter n, size of the array. Each cell PX[i] has 3 components: PX[i] = \{PX[i].ID, PX[i].x, PX[i].y \}: the ID, $x$ and $y$ coordinate of point PX[i]
The points in PX are sorted (ascending order) according to the x coordinate
//Output: min_dist: minimum distant between two points in PX
PY[0 .. n-1]: Points in PX now sorted (ascending order) according to the y coordinate
//check the base case

If $\mathrm{n}=1$ return (INF, PX)
If $\mathrm{n}=2$
min_dist $\leftarrow \operatorname{dist}(\operatorname{PX}[0], \mathrm{PX}[1])$
If PX[0]. $\mathrm{y} \leq \mathrm{PX}[1] . \mathrm{y}$
$\mathrm{PY}[0] \leftarrow \mathrm{PX}[0]$
$\mathrm{PY}[1] \leftarrow \mathrm{PY}[1]$
Else PY[0] $\leftarrow \mathrm{PX}[1]$
$\mathrm{PY}[1] \leftarrow \mathrm{PX}[0]$
Return (min_dist, PY)
//Divide into two subproblems
$\operatorname{mid} \leftarrow \mathrm{n} / 2$
PXL $\leftarrow$ PX[0..mid-1]
PXR $\leftarrow$ PX[mid .. $\mathrm{n}-1]$
// Conquer the subproblems
(dL, PYL) $\leftarrow$ Closest_pair $($ PXL, mid) $\quad / /$ the left half of the points // PYL is the array of points in PXL sorted by y coordinate
$(\mathrm{dR}, \mathrm{PYR}) \leftarrow$ Closest_pair $(\mathrm{PXR}, \mathrm{n}-\mathrm{mid}) \quad / /$ the right half of the points
// PYR is the array of points in PXR sorted by y coordinate
//Combine the solutions for the subproblems
$\mathrm{d} \leftarrow \min (\mathrm{dL}, \mathrm{dR}) \quad / / \mathrm{d}$ is the current minimum distance
$\mathrm{PY} \leftarrow$ merge $(\mathrm{PYL}, \mathrm{PYR}) \quad / / \mathrm{PY}$ is the array of points in PX sorted by the y coordinate
// The merge method is essentially the same as in merge sort

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mid_x }\leftarrow\mathrm{ PX[mid].x // the x value of the split point
length }\leftarrow
i}\leftarrow
While (i \leqn-1) Do
    If |PY[i].x - mid_x | \leq d
        Strip[length] \leftarrowPY[i] // point PY[i] is within the strip of width 2d centered around
                // the line x = mid_x
            length }\leftarrow\mathrm{ length +1
    i}\leftarrow\textrm{i}+
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                // Now the array Strip contains all points in PY which are within
                    // the strip of width 2d centered around the line \(x=\) mid_x
                    //The number of elements in Strip is length
    min_dist $\leftarrow \mathrm{d}$
//Next we will check the points in Strip for possible smaller
// distance than min_dist

For $\mathrm{i} \leftarrow 0$ to length-2 Do
$\mathrm{k} \leftarrow \mathrm{i}+1 \quad / /$ Only check points in Strip with index larger than i
While ( $k \leq$ length-1 AND Strip[k].y - Strip[i].y $\leq \mathrm{d}$ ) Do
new_d $\leftarrow \operatorname{dist}($ Strip[i], Strip[k])
min_dist $\leftarrow$ min(d, new_d) // update the current min_dist
Return (min_dist, PY)

